Introduction to OpenMP

March 29, 2017



Setup Examples

To copy the demos and examples to your local scratch, type:

<u>/general/public/training/openmp/setup.sh</u> (ADA & CURIE) <u>/sw/local/training/openmp/setup.sh</u> (TERRA)

Basic Computer Architecture

Each (terra/ada/curie) **NODE** contains multiple cores (28 on terra, 20 on ada, 16 on curie) and at at least 64GB (256 on curie) of **Shared** memory (i.e. all cores have access to this memory)



NOTE: we use ada architecture for illustration below

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What is OpenMP?

Defacto standard API for writing <u>shared memory</u> parallel applications in C, C++, and Fortran

OpenMP API consists of:

- Compiler pragmas/directives
- Runtime subroutines/functions
- Environment variables

OpenMP Directives/Pragmas

fortran directive format:

!\$OMP DIRECTIVE [clauses]
.
.
!\$OMP END DIRECTIVE



C/C++ pragma format is:



All OpenMP directives follow this format.

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OpenMP HelloWorld



COMPILING:

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intel: ifort -openmp -o hi.x hello.f90intel: icpc -openmp -o hi.x hello.cpppgi: pgfortran -mp -o hi.x hello.f90pgi: pgcpp -mp -o hi.x hello.cppgnu: gfortran -fopenmp -o hi.x hello.f90gnu: g++ -fopenmp -o hi.x hello.cppIBM: xlf_r -qsmp=omp -o hi.x hello.f90IBM: xlc++_r -qsmp=omp -o hi.x hello.cpp



Fork/Join



OpenMP follows the fork/join model:

- OpenMP programs start with a single thread; the master thread (Thread #0)
- At start of parallel region master starts team of parallel "worker" threads (FORK)
- Statements in parallel block are executed in parallel by every thread
- At end of parallel region, all threads synchronize, and join master thread (JOIN)

Implicit barrier. Will discuss synchronization later

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Cores & Threads

What are threads, cores, and how do they relate?

- > Thread is independent sequence of execution of program code
 - Block of code with one entry and one exit
- For our purposes a more abstract concept
- Unrelated to physical cores/CPUs
- OpenMP threads are mapped onto physical cores
- Possible to map more than 1 thread onto a core
- In practice best to have one-to-one mapping.

Setting the number of Threads

Environmental variable: OMP_NUM_THREADS

```
export OMP_NUM_THREADS=4 ./a.out
```

Runtime function: omp_set_num_threads(n)

```
omp_set_num_threads(4);
!$omp parallel
```

OMP PARALLEL clause: num_threads(n)

```
!$omp parallel num_threads(4)
```

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case sensitive

???

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Useful OpenMP Functions

- Runtime function omp_get_num_threads()
 - Returns number of threads in parallel region
 - ✓ Returns 1 if called outside parallel region
- Runtime function omp_get_thread_num()
 - Returns id of thread in team
 - ✓ Value between [0,n-1] // where n = #threads
 - ✓ Master thread always has id 0
- Runtime function omp_get_max_threads()
 - Returns upper bound #threads in parallel region
 - value of OMP_NUM_THREADS or set_num_threads

DEMO 1

We will create an OpenMP program that does the following:

- 1) print total number of threads
- 2) start parallel region
- 3) every thread prints it's id and total number of threads
- 4) close the parallel region

Remember: memory is (conceptually) shared by all threads

\$OMP PARALLEL
TOT = omp_get_num_threads()
ID = omp_get_thread_id()

!\$OMP END PARALLEL

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Remember: memory is (conceptually) shared by all threads

\$OMP PARALLEL
TOT = omp_get_num_threads()
ID = omp_get_thread_id()

!\$OMP END PARALLEL

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All threads try to access the same variable (possibly at the same time). This can lead to a <u>race</u> condition. Different runs of same program might give different results because of these race conditions



Data Scope Clauses

Data scope clauses: private(list)

!\$OMP PARALLEL PRIVATE(a,c)
 :
!\$OMP END PARALLEL

#pragma omp parallel private(a,c)

- Every thread will have it's own "private" copy of variables in list
- > No other thread has access to this "private" copy
- Private variables are NOT initialized with value before region started.
- Private variables are NOT accessible after enclosing region finishes

Index variables (Fortran, C/C++) and variables declared inside parallel region (C/C++) are considered private by default.

Data Scope Clauses

Data scope clauses: shared(list)

!\$OMP PARALLEL SHARED(a,c)
 :
!\$OMP END PARALLEL

#pragma omp parallel shared(a,c)

- All variables in list will be considered shared.
- Every openmp thread has access to all these variables
- Programmer's responsibility to avoid race conditions

By default most variables in work sharing constructs are considered shared in OpenMP. Exceptions include index variables (Fortran, C/C++) and variables declared inside parallel region (C/C++).

Data Scope Clauses

Data scope clauses: firstprivate(list)

!\$OMP PARALLEL FIRSTPRIVATE(a,c)
!\$OMP END PARALLEL

#pragma omp parallel firstprivate(a,c)

- Very similar to private clause
- Private copies are initialized with value of original variable

Data scope clauses: default(shared | private | firstprivate | lastprivate)

!\$OMP PARALLEL DEFAULT(private)
!\$OMP END PARALLEL

#pragma omp parallel default(private)

Sets default data scoping rule

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- If not set, default depends on directive
 - e.g. shared for work sharing directives

DEMO 2

We will rewrite the OpenMP program (from previous demo) and make sure all variables are assigned and printed correctly.

Remember: memory is (conceptually) shared by all threads

\$OMP PARALLEL PRIVATE(ID)
TOT = omp_get_num_threads()
ID = omp_get_thread_id()

!\$OMP END PARALLEL

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Private memories

TIP 1: Stack size

- OpenMP creates separate data stack for every worker thread to store copies of private variables (master thread uses regular stack)
 Size of these stacks is not defined by OpenMP standards
 Behavior of program undefined when stack space exceeded
 Although most compilers/RT will throw seg fault
 To increase stack size use environment var OMP_STACKSIZE, e.g.
 export OMP_STACKSIZE=512M
 - ✓ export OMP_STACKSIZE=1G
- To make sure master thread has large enough stack space use ulimit -s command (unix/linux).

Work Sharing Directives (FORTRAN)

Work sharing directive (Fortran): **!\$OMP DO** [clauses]

!\$OMP PARALLEL
!\$OMP DO
DO n=1,N
A(n) = A(n) + B
ENDDO
!\$OMP END DO
!\$OMP END PARALLEL

!\$OMP PARALLEL DO
DO n=1,N
A(n) = A(n) + B
ENDDO
!\$OMP END PARALLEL DO

DO command must immediately follow "!\$OMP DO" directive

- Loop iteration variable is "private" by default
- If "end do" directive omitted it is assumed at end of loop

OR

Not case sensitive

Work Sharing Directives (C/C++)

Work sharing pragma (C/C++): **#pragma omp for** [clauses]

#pragma omp parallel
#pragma omp for
for (int i=1;i<N;++i)
 A(n) = A(n) + B;</pre>

OR

#pragma omp parallel for for (int i=1;i<N;++i) A(n) = A(n) + B;

- for command must immediately follow "#pragma omp for"
- Newline required after "#pragma omp for"
- Iteration variable can be
 - ✓ Signed/unsigned integer variable
 - Pointer type
 - ✓ Random access iterator

Work Sharing Directives (C/C++)

C++

new in OpenMP 3.0.

Random access iterators:

vector<int> vec(10); vector<int>::iterator it= vec.begin(); #pragma omp parallel for for (; it != vec.end() ; ++it) { // do something with *it }

Pointer type:

int N = 1000000; int arr[N]; #pragma omp parallel for for (int* t=arr;t<arr+N;++t) { // do something with *t }

DEMO 3

We will create a program that computes a simple matrix vector multiplication b=Ax, either in fortran or C/C++. We will use OpenMP directives (pragmas) to make it run in parallel.

NOTE: demo matrix_mult

Data Dependencies

Not all loops can be parallelized. Before adding OpenMP directives we need to check for <u>data dependencies</u>

For example:

Is the result guranteed to be correct if you run this loop in parallel?

Data Dependencies

Not all loops can be parallelized. Before adding OpenMP directives we need to check for <u>data dependencies</u>

For example:

Easier to see when you unroll loop (partly):

iteration i=1:

iteration i=2:

iteration i=3:

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REDUCTION

Data scope clause: REDUCTION(op:list)

In a reduction the values of the local copies will be summarized (reduced) into a global shared variable (using the specified reduction operator).

Example:

for (int i=0;i<10;++i) sum=sum+a[i]; #pragma omp parallel for reduction(+:sum)
for (int i=0;i<10;++i)
 sum=sum+a[i];</pre>

Only certain kind of operators allowed

 +, -, *, max, min______
 &, |, ^, &&, || (C++)
 Also allowed in c/c++ since OpenMP 3.1
 .and., .or., .eqv., .neqv., iand, ior, ieor (Fortran)
 OpenMP 4.0 allows for user defined reductions

 Variables in list have to be *shared*

DEMO 4

We will create a subroutine that computes the dot product of two vectors. We will use OpenMP pragmas (directives) to make it run in parallel.

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TIP 2: ORPHANED DIRECTIVES

An OpenMP directive (pragma) that appears independently from another enclosing directive is called an <u>orphaned directive</u> (pragma). It exists outside of another directive's (pragma) static (lexical) extent. For example:

int main() {
#pragma omp parallel
foo()
return 0;
}

void foo() {
#pragma omp for
for (int i=0;i<N;i++) {....}
}</pre>

Note: OpenMP directives (pragmas) should be in the dynamic extent of a parallel section directive (pragma).

SCHEDULE (STATIC,250) //loop with 1000 iterations, 4 threads

!\$OMP PARALLEL DO SCHEDULE (STATIC,250)
DO i=1,1000
...
ENDDO
!\$OMP END PARALLEL DO



Although the OpenMP standard does not specify how a loop should be partitioned most compilers split the loop in N/p (N #iterations, p #threads) chunks by default.

SCHEDULE (STATIC,10) //loop with 1000 iterations, 4 threads







With static scheduling the number of iterations is evenly distributed among all openmp threads (i.e. Every thread will be assigned similar number of iterations). This is not always the best way to partition. Why is This?



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SCHEDULE (DYNAMIC,10) //loop with 1000 iterations, 4 threads

!\$OMP PARALLEL DO SCHEDULE (DYNAMIC,10)
DO i=1,1000
...
ENDDO
!\$OMP END PARALLEL DO

#pragma omp parallel for schedule(dynamic,10)
for (int i=0;i<1000;++i) {
 :
 }</pre>

Loop iterations are divided into pieces of size chunk. When a thread finishes one chunk, it is dynamically assigned another.

NOTE: there is a significant overhead involved compared to static scheduling. WHY?

SCHEDULE (GUIDED,10) //loop with 1000 iterations, 4 threads

!\$OMP PARALLEL DO SCHEDULE (GUIDED,10)
DO i=1,1000
...
ENDDO
!\$OMP END PARALLEL DO

Similar to DYNAMIC schedule except that chunk size is relative to number of iterations left.

NOTE: there is a significant overhead involved compared to static scheduling. WHY?

Work Sharing Directives (2)

#pragma omp single (!\$OMP SINGLE)

Although technically it doesn't share work, another work sharing pragma that OpenMP provides is **#pragma omp single (!\$OMP SINGLE)** When encountering a single directive only one member of the team will execute the code in the block

- One thread (not neccesarily master) executes the block
- Other threads will wait
- Useful for thread-unsafe code
- Useful for I/O operations

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Work Sharing Directives (3)

#pragma omp sections (!\$OMP SECTIONS)

Executes set of structured blocks in parallel

#pragma omp parallel
#pragma omp sections
{
#pragma omp section
 // WORK 1
#pragma omp section
 // WORK 2
}

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!\$OMP PARALLEL
!\$OMP SECTIONS

!\$OMP SECTION
 // WORK 1
!\$OMP SECTION
 // WORK 2
!\$OMP END SECTIONS
!\$OMP END PARALLEL

This will execute "WORK 1" and "WORK 2" in parallel

NOWAIT Clause

Worksharing constructs have an implicit barrier at the end of their worksharing region. However, OpenMP provides a clause to ommit this barrier.

#pragma omp for nowait

!\$OMP DO : !\$OMP END DO <u>NOWAIT</u>

At end of work sharing constructs threads will not wait
 There is always barrier at end of parallel region

Work Sharing Summary

OMP work sharing constructions discussed

- #pragma omp for
- #pragma omp sections
- #pragma omp single

(!\$OMP DO) (!\$OMP SECTIONS) (!\$OMP SINGLE)

Clauses that can be used with these constructs (

incomplete list and not all clauses can be used with every directive)

- SHARED (list)
- PRIVATE (list)
- ➢ FIRSTPRIVATE (list)
- LASTPRIVATE(list)
- SCHEDULE (STATIC | DYNAMIC | GUIDED, chunk)
- REDUCTION(op:list)
- NOWAIT

Communication/Synchronization

OpenMP programs use shared variables to communicate. We need to make sure these variables are not accessed at the same time by different threads (will cause race conditions). OpenMP provides a number of directives for synchronization.

Synchronization Directives (1)

#pragma omp master (!\$OMP MASTER)

This Directive ensures that only the master threads excecutes instructions in the block. There is no implicit barrier so other threads will not wait for master to finish

What is difference with **!**\$OMP SINGLE DIRECTIVE?

Synchronization Directives (2)

#pragma omp critical (!\$OMP CRITICAL)

This Directive makes sure that only one thread can execute the code in the block. If another threads reaches the critical section it will wait untill the current thread finishes this critical section. <u>Every thread</u> will execute the critical block and they will synchronize at end of critical section

Introduces overhead

- Serializes critical block
- \blacktriangleright If time in critical block relatively large \rightarrow speedup negliable

DEMO 5

In the REDUCTION exercise we created an OpenMP program that computes the dotproduct of two vectors using the reduction clause. Now we will create a version that uses critical blocks instead.

NOTE: dempo dotproduct_critical

Synchronization Directives (3)

#pragma omp atomic (!\$OMP ATOMIC)

This Directive is very similar to the **!**\$OMP CRITICAL directive on the previous slide. Difference is that **!**\$OMP ATOMIC is only used for the update of a memory location. Sometimes **!**\$OMP ATOMIC is also referred to as a mini critical section.

- Block consists of only one statement
- Atomic statement must follow specific syntax

Synchronization Directives (4)

#pragma omp barrier (!\$OMP BARRIER)

A barrier will force every thread to wait at the barrier until all threads have reached the barrier. The following omp directives we discussed before include an implicit barrier:

- !\$ OMP END PARALLEL
- > !\$ OMP END DO
- IS OMP END SECTIONS
- IS OMP END SINGLE
- IS OMP END CRITICAL

TIP 3: IF Clause

OpenMP provides another useful clause to decide at run time if a parallel region should actually be run in parallel (multiple threads) or just by the master thread:

IF (logical expr)

For example:

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\$!OMP PARALLEL IF(n > 100000) #pragma omp parallel if (n>100000) (fortran) (C/C++)

This will only run the parallel region when n > 100000

Nested Parallelism

Nested Parallelism

OpenMP allows parallel regions inside other parallel regions ,e.g.

#pragma omp parallel for

```
for (int i=0; i<N;++i) {
```

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```
#pragma omp parallel for
for (j=0;j<M;++j)
}
```

- > To enable nested parallelism:
 - ✓ env var: OMP_NESTED=1
 - ✓ lib function: omp_set_nested(1)
- To specify number of threads:
 - ✓ omp_set_ num_threads()
 - ✓ OMP_NUM_THREADS=4,2

NOTE: using nested parallelism does introduce extra overhead and there is a possibility of over-subscription of threads

MKL

The Intel Math Kernel Library (MKL) has very specialized and optimized versions of many math functions (e.g. blas, lapack). Many of these have been parallelized using OpenMP.

To run mkl functions in parallel use:

- MKL_NUM_THREADS
- OMP_NUM_THREADS

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For more information about MKL:

http://hprc.tamu.edu/wiki/index.php/Ada:MKL